

Dedicated to the memory of Academician Mitrofan M. Cioban (1942-2021)

Boolean asynchronous systems vs. Daizhan Cheng's theory

SERBAN E. VLAD

Abstract. The theory of Daizhan Cheng [1] replaces $\mathbf{B} = \{0, 1\}$ with $\mathbf{D} = \left\{ \begin{pmatrix} 1 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 1 \end{pmatrix} \right\}$, and Boolean functions with logical matrices. Interesting and very important algebraical opportunities result, which can be used in systems theory. Our purpose is to give a hint on the theory of Cheng and its application to asynchronicity.

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Sisteme booleene asincrone din perspectiva teoriei lui Daizhan Cheng

Rezumat. Teoria lui Daizhan Cheng [1] înlocuiește $\mathbf{B} = \{0, 1\}$ cu $\mathbf{D} = \left\{ \begin{pmatrix} 1 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 1 \end{pmatrix} \right\}$, și funcțiile booleene cu matrici logice. Rezultă de aici oportunități algebrice importante, care pot fi folosite în teoria sistemelor. Scopul nostru este acela de a schița teoria lui Cheng și aplicațiile sale în asincronism.

Cuvinte cheie: funcție booleană, sistem asincron boolean, matrice de structură, produs semi-tensorial, teoria lui Daizhan Cheng.

1. PRELIMINARIES

Notation 1.1. We denote with $\mathbf{B} = \{0, 1\}$ the binary Boolean algebra.

Definition 1.1. The λ -iterate of $\Phi : \mathbf{B}^n \rightarrow \mathbf{B}^n$, $\lambda \in \mathbf{B}^n$ is the function $\Phi^\lambda : \mathbf{B}^n \rightarrow \mathbf{B}^n$ defined like this: $\forall \mu \in \mathbf{B}^n, \forall i \in \{1, \dots, n\}$,

$$(\Phi^\lambda)_i(\mu) = \begin{cases} \Phi_i(\mu), & \text{if } \lambda_i = 1, \\ \mu_i, & \text{if } \lambda_i = 0. \end{cases}$$

Dedicated to the memory of Academician Mitrofan Cioban, who was the core of the intellectual and spiritual life of the Moldovan mathematicians for so many years, and also a steady bridge connecting the mathematicians from Moldova and Romania. We shall keep in our hearts his common sense and support. May he rest in peace!

Definition 1.2. Given Φ , the function $\tilde{\Phi} : \mathbf{B}^n \times \mathbf{B}^n \rightarrow \mathbf{B}^n$ is defined by $\forall \mu \in \mathbf{B}^n, \forall \lambda \in \mathbf{B}^n$,

$$\tilde{\Phi}(\mu, \lambda) = \Phi^\lambda(\mu). \quad (1)$$

Definition 1.3. The function $\alpha : \mathbf{N} \rightarrow \mathbf{B}^n, \mathbf{N} \ni k \mapsto \alpha^k \in \mathbf{B}^n$, with

$$\forall i \in \{1, \dots, n\}, \text{ the set } \{k \mid k \in \mathbf{N}, \alpha_i^k = 1\} \text{ is infinite}$$

is called *progressive computation function*, and we denote with Π_n the set of these functions.

Remark 1.1. Two ways of making the discrete time iterations of the function $\Phi : \mathbf{B}^n \rightarrow \mathbf{B}^n$ exist: synchronously $1_{\mathbf{B}^n}, \Phi, \Phi \circ \Phi, \dots$ when Φ_1, \dots, Φ_n are computed always, all of them, and asynchronously, when the coordinates of Φ are computed sometimes, independently on each other. The functions $\alpha \in \Pi_n$ indicate how Φ is computed: $\forall k \in \mathbf{N}, \forall i \in \{1, \dots, n\}$,

$$\begin{cases} \alpha_i^k = 1, & \text{at time instant } k, \Phi_i \text{ is computed,} \\ \alpha_i^k = 0, & \text{at time instant } k, \Phi_i \text{ is not computed.} \end{cases}$$

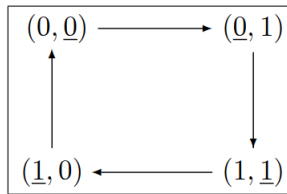
Synchronicity is that special case of asynchronicity when $\forall k \in \mathbf{N}, \alpha^k = (1, \dots, 1)$.

Definition 1.4. The *unbounded delay model of computation of Φ* consists in the equation

$$x(k+1) = \Phi^{\alpha^k}(x(k)), \quad (2)$$

where $\Phi : \mathbf{B}^n \rightarrow \mathbf{B}^n, x : \mathbf{N} \rightarrow \mathbf{B}^n, \alpha \in \Pi_n$ and $k \in \mathbf{N}$. In (2) the function x , called *state*, is unknown, and $x(0)$, together with α , are parameters.

Example 1.1. We consider the function $\Phi : \mathbf{B}^2 \rightarrow \mathbf{B}^2, \forall \mu \in \mathbf{B}^2, \Phi(\mu_1, \mu_2) = (\mu_2, \overline{\mu_1})$, with the following state portrait



In the drawing, the underlined coordinates $\underline{\mu}_i, i \in \{1, 2\}$ show that $\Phi_i(\mu) \neq \mu_i$ and, by their computation, the system moves to a distinct state, while the arrows indicate the evolution of the system. The equation (2) is $\forall k \in \mathbf{N}$,

$$\begin{cases} x_1(k+1) = x_2(k)\alpha_1^k \cup x_1(k)\overline{\alpha_1^k}, \\ x_2(k+1) = \overline{x_1(k)}\alpha_2^k \cup x_2(k)\overline{\alpha_2^k}, \end{cases} \quad (3)$$

where $x : \mathbf{N} \rightarrow \mathbf{B}^2$ fulfils $x(0) = (0, 0)$ and $\alpha \in \Pi_2$ is defined as

$$\alpha = (1, 0), (0, 1), (1, 1), (0, 1), (1, 0), \dots$$

We get

$$x(1) = \Phi^{\alpha^0}(x(0)) = \Phi^{(1,0)}(0, 0) = (0, 0), \quad (4)$$

$$x(2) = \Phi^{\alpha^1}(x(1)) = \Phi^{(0,1)}(0, 0) = (0, 1), \quad (5)$$

$$x(3) = \Phi^{\alpha^2}(x(2)) = \Phi^{(1,1)}(0, 1) = (1, 1), \quad (6)$$

$$x(4) = \Phi^{\alpha^3}(x(3)) = \Phi^{(0,1)}(1, 1) = (1, 0), \quad (7)$$

$$x(5) = \Phi^{\alpha^4}(x(4)) = \Phi^{(1,0)}(1, 0) = (0, 0), \quad (8)$$

...

2. SEMI-TENSOR PRODUCT

Notation 2.1. We use the notation $M_{m \times n}$ for the set of the matrices with binary entries that have m rows and n columns.

Remark 2.1. In the following Definitions 2.1 and 2.2, the operations with matrices are induced by the field structure of \mathbf{B} relative to \oplus, \cdot .

Definition 2.1. The *Kronecker product* \otimes of the matrices $A \in M_{m \times n}$ and $B \in M_{p \times q}$ is

$$A \otimes B = \begin{pmatrix} a_{11}B & \dots & a_{1n}B \\ \dots & \dots & \dots \\ a_{m1}B & \dots & a_{mn}B \end{pmatrix} \in M_{mp \times nq}.$$

Definition 2.2. The *semi-tensor product* \ltimes of $A \in M_{m \times n}$ and $B \in M_{p \times q}$ is by definition

$$A \ltimes B = (A \otimes I_n^c)(B \otimes I_p^c) \in M_{\frac{mc}{n} \times \frac{qc}{p}},$$

where I_k is the $k \times k$ identity matrix and c is the least common multiple of n and p .

Remark 2.2. At Definition 2.2, $A \otimes I_n^c$ has $n \frac{c}{n}$ columns and $B \otimes I_p^c$ has $p \frac{c}{p}$ rows, thus the product of the matrices $A \otimes I_n^c, B \otimes I_p^c$ makes sense.

Remark 2.3. If $n = p$, the semi-tensor product coincides with the usual product of the matrices. This happens because we get $c = n = p$, $A \otimes I_1 = A$, and $B \otimes I_1 = B$.

Example 2.1. We have the following examples of Kronecker product

$$\begin{pmatrix} 1 \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 1 & 1 \end{pmatrix} = \begin{pmatrix} 1 \begin{pmatrix} 1 & 1 \end{pmatrix} \\ 0 \begin{pmatrix} 1 & 1 \end{pmatrix} \end{pmatrix} = \begin{pmatrix} 1 & 1 \\ 0 & 0 \end{pmatrix},$$

and semi-tensor product

$$\begin{pmatrix} 1 \\ 0 \end{pmatrix} \ltimes \begin{pmatrix} 1 & 1 & 1 \\ 0 & 0 & 0 \end{pmatrix} = \left(\begin{pmatrix} 1 \\ 0 \end{pmatrix} \otimes I_2 \right) \left(\begin{pmatrix} 1 & 1 & 1 \\ 0 & 0 & 0 \end{pmatrix} \otimes I_1 \right)$$

$$= \begin{pmatrix} 1 & 0 \\ 0 & 1 \\ 0 & 0 \\ 0 & 0 \end{pmatrix} \begin{pmatrix} 1 & 1 & 1 \\ 0 & 0 & 0 \end{pmatrix} = \begin{pmatrix} 1 & 1 & 1 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix}.$$

Remark 2.4. The semi-tensor product is associative, and for this reason we shall omit writing brackets when it is used repeatedly.

3. REPLACEMENT OF \mathbf{B} WITH \mathbf{D}

Notation 3.1. We denote with $\delta_n^i \in M_{n \times 1}$ the columns of the identity matrix of dimension n :

$$\delta_n^i = \begin{pmatrix} 0 \\ \dots \\ 1 \\ \dots \\ 0 \end{pmatrix} - i,$$

where $n \geq 1$ and $i \in \{1, \dots, n\}$.

Notation 3.2. We use also the notations

$$\mathbf{D} = \{\delta_2^1, \delta_2^2\},$$

$$\mathbf{D}^{(n)} = \{\delta_{2^n}^1, \dots, \delta_{2^n}^{2^n}\}.$$

Remark 3.1. \mathbf{D} and $\mathbf{D}^{(n)}$ do not have a name and an algebraical structure of their own, but they will act as \mathbf{B} and \mathbf{B}^n in the following. Obviously, $card(\mathbf{B}) = card(\mathbf{D}) = 2$ and $card(\mathbf{B}^n) = card(\mathbf{D}^{(n)}) = 2^n$.

Notation 3.3. We use the notations $\zeta : \mathbf{B} \rightarrow \mathbf{D}$, $\zeta_n : \mathbf{B}^n \rightarrow \mathbf{D}^{(n)}$ for the following functions: $\forall \mu \in \mathbf{B}, \forall \lambda \in \mathbf{B}^n$,

$$\zeta(\mu) = \begin{pmatrix} \mu \\ \bar{\mu} \end{pmatrix},$$

$$\zeta_n(\lambda) = \begin{pmatrix} \lambda_1 \dots \lambda_{n-1} \lambda_n \\ \lambda_1 \dots \lambda_{n-1} \bar{\lambda}_n \\ \lambda_1 \dots \bar{\lambda}_{n-1} \lambda_n \\ \dots \\ \bar{\lambda}_1 \dots \bar{\lambda}_{n-1} \bar{\lambda}_n \end{pmatrix}.$$

We denote in general $\underline{\underline{\mu}} = \zeta(\mu)$ and $\underline{\underline{\lambda}} = \zeta_n(\lambda)$.

Remark 3.2. We notice that for any $\mu \in \mathbf{B}$, respectively $\lambda \in \mathbf{B}^n$, exactly one of $\mu, \bar{\mu}$ is 1, respectively exactly one of $\lambda_1 \dots \lambda_{n-1} \lambda_n, \lambda_1 \dots \lambda_{n-1} \bar{\lambda}_n, \lambda_1 \dots \bar{\lambda}_{n-1} \lambda_n, \dots, \bar{\lambda}_1 \dots \bar{\lambda}_{n-1} \bar{\lambda}_n$ is 1, meaning that $\underline{\underline{\mu}} \in \mathbf{D}$, respectively that $\underline{\underline{\lambda}} \in \mathbf{D}^{(n)}$ indeed.

Theorem 3.4. (a) ζ and ζ_n are bijections;

(b) $\forall \lambda \in \mathbf{B}^n$,

$$\underline{\underline{\lambda}} = \underline{\underline{\lambda_1}} \times \dots \times \underline{\underline{\lambda_n}}.$$

Proof. (a) When $\lambda \in \mathbf{B}^n$ takes the distinct 2^n values $(1, \dots, 1, 1), (1, \dots, 1, 0), (1, \dots, 0, 1), \dots, (0, \dots, 0, 0)$, $\underline{\underline{\lambda}}$ takes the distinct 2^n values $\delta_{2^n}^1, \delta_{2^n}^2, \delta_{2^n}^3, \dots, \delta_{2^n}^{2^n}$.

(b) For $n = 2$ and arbitrary $\lambda \in \mathbf{B}^2$, we obtain

$$\begin{aligned} \underline{\underline{\lambda_1}} \times \underline{\underline{\lambda_2}} &= \begin{pmatrix} \lambda_1 \\ \bar{\lambda}_1 \end{pmatrix} \times \begin{pmatrix} \lambda_2 \\ \bar{\lambda}_2 \end{pmatrix} = \left(\begin{pmatrix} \lambda_1 \\ \bar{\lambda}_1 \end{pmatrix} \otimes \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \right) \left(\begin{pmatrix} \lambda_2 \\ \bar{\lambda}_2 \end{pmatrix} \otimes 1 \right) \\ &= \begin{pmatrix} \lambda_1 & 0 \\ 0 & \lambda_1 \\ \bar{\lambda}_1 & 0 \\ 0 & \bar{\lambda}_1 \end{pmatrix} \begin{pmatrix} \lambda_2 \\ \bar{\lambda}_2 \end{pmatrix} = \begin{pmatrix} \lambda_1 \lambda_2 \\ \lambda_1 \bar{\lambda}_2 \\ \bar{\lambda}_1 \lambda_2 \\ \bar{\lambda}_1 \bar{\lambda}_2 \end{pmatrix} = \underline{\underline{\lambda}}. \end{aligned}$$

The property is supposed to be true for n and the proof is made for $n + 1$. □

4. STRUCTURE MATRIX

Notation 4.1. The notation of the i -th column of an arbitrary binary matrix A is $\text{col}_i(A)$.

Definition 4.1. A matrix A with n rows and m columns is called *logical* if $\forall j \in \{1, \dots, m\}, \text{col}_j(A) \in \{\delta_n^1, \dots, \delta_n^n\}$. The set of the logical matrices with n rows and m columns is denoted with $L_{n \times m}$.

Definition 4.2. Let $f : \mathbf{B}^n \rightarrow \mathbf{B}$, $\Phi : \mathbf{B}^n \rightarrow \mathbf{B}^n$ and $\tilde{\Phi} : \mathbf{B}^n \times \mathbf{B}^n \rightarrow \mathbf{B}^n$, as defined at (1).

We denote with $M_f \in L_{2 \times 2^n}$ the matrix

$$M_f = \begin{pmatrix} \frac{f(1, \dots, 1, 1)}{f(1, \dots, 1, 1)}, & \frac{f(1, \dots, 1, 0)}{f(1, \dots, 1, 0)}, & \frac{f(1, \dots, 0, 1)}{f(1, \dots, 0, 1)}, & \dots & \frac{f(0, \dots, 0, 0)}{f(0, \dots, 0, 0)} \end{pmatrix},$$

with $M_\Phi \in L_{2^n \times 2^n}$ the matrix whose columns are

$$\text{col}_1(M_\Phi) = \begin{pmatrix} \Phi_1(1, \dots, 1, 1) \dots \Phi_{n-1}(1, \dots, 1, 1) \Phi_n(1, \dots, 1, 1) \\ \Phi_1(1, \dots, 1, 1) \dots \Phi_{n-1}(1, \dots, 1, 1) \bar{\Phi}_n(1, \dots, 1, 1) \\ \Phi_1(1, \dots, 1, 1) \dots \Phi_{n-1}(1, \dots, 1, 1) \Phi_n(1, \dots, 1, 1) \\ \dots \\ \bar{\Phi}_1(1, \dots, 1, 1) \dots \bar{\Phi}_{n-1}(1, \dots, 1, 1) \bar{\Phi}_n(1, \dots, 1, 1) \end{pmatrix},$$

$$\text{col}_2(M_\Phi) = \begin{pmatrix} \Phi_1(1, \dots, 1, 0) \dots \Phi_{n-1}(1, \dots, 1, 0) \overline{\Phi_n(1, \dots, 1, 0)} \\ \Phi_1(1, \dots, 1, 0) \dots \Phi_{n-1}(1, \dots, 1, 0) \Phi_n(1, \dots, 1, 0) \\ \Phi_1(1, \dots, 1, 0) \dots \overline{\Phi_{n-1}(1, \dots, 1, 0)} \overline{\Phi_n(1, \dots, 1, 0)} \\ \dots \\ \overline{\Phi_1(1, \dots, 1, 0)} \dots \overline{\Phi_{n-1}(1, \dots, 1, 0)} \overline{\Phi_n(1, \dots, 1, 0)} \\ \dots \\ \Phi_1(0, \dots, 0, 0) \dots \Phi_{n-1}(0, \dots, 0, 0) \overline{\Phi_n(0, \dots, 0, 0)} \\ \Phi_1(0, \dots, 0, 0) \dots \Phi_{n-1}(0, \dots, 0, 0) \Phi_n(0, \dots, 0, 0) \\ \Phi_1(0, \dots, 0, 0) \dots \overline{\Phi_{n-1}(0, \dots, 0, 0)} \overline{\Phi_n(0, \dots, 0, 0)} \\ \dots \\ \overline{\Phi_1(0, \dots, 0, 0)} \dots \overline{\Phi_{n-1}(0, \dots, 0, 0)} \overline{\Phi_n(0, \dots, 0, 0)} \end{pmatrix},$$

$$\text{col}_{2n}(M_\Phi) = \begin{pmatrix} \Phi_1(0, \dots, 0, 0) \dots \Phi_{n-1}(0, \dots, 0, 0) \overline{\Phi_n(0, \dots, 0, 0)} \\ \Phi_1(0, \dots, 0, 0) \dots \Phi_{n-1}(0, \dots, 0, 0) \Phi_n(0, \dots, 0, 0) \\ \Phi_1(0, \dots, 0, 0) \dots \overline{\Phi_{n-1}(0, \dots, 0, 0)} \overline{\Phi_n(0, \dots, 0, 0)} \\ \dots \\ \overline{\Phi_1(0, \dots, 0, 0)} \dots \overline{\Phi_{n-1}(0, \dots, 0, 0)} \overline{\Phi_n(0, \dots, 0, 0)} \end{pmatrix}$$

and with $M_{\overline{\Phi}} \in L_{2n \times 2n}$ the matrix

$$\text{col}_1(M_{\overline{\Phi}}) = \begin{pmatrix} \Phi_1^{(1, \dots, 1, 1)}(1, \dots, 1) \dots \Phi_{n-1}^{(1, \dots, 1, 1)}(1, \dots, 1) \overline{\Phi_n^{(1, \dots, 1, 1)}(1, \dots, 1)} \\ \Phi_1^{(1, \dots, 1, 1)}(1, \dots, 1) \dots \overline{\Phi_{n-1}^{(1, \dots, 1, 1)}(1, \dots, 1)} \overline{\Phi_n^{(1, \dots, 1, 1)}(1, \dots, 1)} \\ \Phi_1^{(1, \dots, 1, 1)}(1, \dots, 1) \dots \Phi_{n-1}^{(1, \dots, 1, 1)}(1, \dots, 1) \Phi_n^{(1, \dots, 1, 1)}(1, \dots, 1) \\ \dots \\ \overline{\Phi_1^{(1, \dots, 1, 1)}(1, \dots, 1)} \dots \overline{\Phi_{n-1}^{(1, \dots, 1, 1)}(1, \dots, 1)} \overline{\Phi_n^{(1, \dots, 1, 1)}(1, \dots, 1)} \end{pmatrix},$$

$$\text{col}_2(M_{\overline{\Phi}}) = \begin{pmatrix} \Phi_1^{(1, \dots, 1, 0)}(1, \dots, 1) \dots \Phi_{n-1}^{(1, \dots, 1, 0)}(1, \dots, 1) \overline{\Phi_n^{(1, \dots, 1, 0)}(1, \dots, 1)} \\ \Phi_1^{(1, \dots, 1, 0)}(1, \dots, 1) \dots \overline{\Phi_{n-1}^{(1, \dots, 1, 0)}(1, \dots, 1)} \overline{\Phi_n^{(1, \dots, 1, 0)}(1, \dots, 1)} \\ \Phi_1^{(1, \dots, 1, 0)}(1, \dots, 1) \dots \Phi_{n-1}^{(1, \dots, 1, 0)}(1, \dots, 1) \Phi_n^{(1, \dots, 1, 0)}(1, \dots, 1) \\ \dots \\ \overline{\Phi_1^{(1, \dots, 1, 0)}(1, \dots, 1)} \dots \overline{\Phi_{n-1}^{(1, \dots, 1, 0)}(1, \dots, 1)} \overline{\Phi_n^{(1, \dots, 1, 0)}(1, \dots, 1)} \end{pmatrix},$$

$$\text{col}_{2n}(M_{\overline{\Phi}}) = \begin{pmatrix} \Phi_1^{(0, \dots, 0, 0)}(1, \dots, 1) \dots \Phi_{n-1}^{(0, \dots, 0, 0)}(1, \dots, 1) \overline{\Phi_n^{(0, \dots, 0, 0)}(1, \dots, 1)} \\ \Phi_1^{(0, \dots, 0, 0)}(1, \dots, 1) \dots \overline{\Phi_{n-1}^{(0, \dots, 0, 0)}(1, \dots, 1)} \overline{\Phi_n^{(0, \dots, 0, 0)}(1, \dots, 1)} \\ \Phi_1^{(0, \dots, 0, 0)}(1, \dots, 1) \dots \Phi_{n-1}^{(0, \dots, 0, 0)}(1, \dots, 1) \Phi_n^{(0, \dots, 0, 0)}(1, \dots, 1) \\ \dots \\ \overline{\Phi_1^{(0, \dots, 0, 0)}(1, \dots, 1)} \dots \overline{\Phi_{n-1}^{(0, \dots, 0, 0)}(1, \dots, 1)} \overline{\Phi_n^{(0, \dots, 0, 0)}(1, \dots, 1)} \end{pmatrix},$$

$$\text{col}_{2n}(M_{\overline{\Phi}}) = \begin{pmatrix} \Phi_1^{(0, \dots, 0, 0)}(0, \dots, 0) \dots \Phi_{n-1}^{(0, \dots, 0, 0)}(0, \dots, 0) \overline{\Phi_n^{(0, \dots, 0, 0)}(0, \dots, 0)} \\ \Phi_1^{(0, \dots, 0, 0)}(0, \dots, 0) \dots \overline{\Phi_{n-1}^{(0, \dots, 0, 0)}(0, \dots, 0)} \overline{\Phi_n^{(0, \dots, 0, 0)}(0, \dots, 0)} \\ \Phi_1^{(0, \dots, 0, 0)}(0, \dots, 0) \dots \Phi_{n-1}^{(0, \dots, 0, 0)}(0, \dots, 0) \Phi_n^{(0, \dots, 0, 0)}(0, \dots, 0) \\ \dots \\ \overline{\Phi_1^{(0, \dots, 0, 0)}(0, \dots, 0)} \dots \overline{\Phi_{n-1}^{(0, \dots, 0, 0)}(0, \dots, 0)} \overline{\Phi_n^{(0, \dots, 0, 0)}(0, \dots, 0)} \end{pmatrix}.$$

$M_f, M_\Phi, M_{\tilde{\Phi}}$ are called the *structure matrices* of $f, \Phi, \tilde{\Phi}$.

Theorem 4.2. *We consider $f : \mathbf{B}^n \rightarrow \mathbf{B}, \Phi : \mathbf{B}^n \rightarrow \mathbf{B}^n$ and $\tilde{\Phi} : \mathbf{B}^n \times \mathbf{B}^n \rightarrow \mathbf{B}^n$ like previously. The assignments*

$$\underline{\underline{\mu}} \mapsto M_f \times \underline{\underline{\mu}},$$

$$\underline{\underline{\mu}} \mapsto M_\Phi \times \underline{\underline{\mu}},$$

$$(\underline{\underline{\mu}}, \underline{\underline{\lambda}}) \mapsto M_{\tilde{\Phi}} \times \underline{\underline{\mu}} \times \underline{\underline{\lambda}},$$

with $\mu \in \mathbf{B}^n, \lambda \in \mathbf{B}^n$, define the functions $M(f) : \mathbf{D}^{(n)} \rightarrow \mathbf{D}, M(\Phi) : \mathbf{D}^{(n)} \rightarrow \mathbf{D}^{(n)}, M(\tilde{\Phi}) : \mathbf{D}^{(n)} \times \mathbf{D}^{(n)} \rightarrow \mathbf{D}^{(n)}$ in the following way: $\forall \underline{\underline{\mu}} \in \mathbf{D}^{(n)}, \forall \underline{\underline{\lambda}} \in \mathbf{D}^{(n)}$,

$$M(f)(\underline{\underline{\mu}}) = M_f \times \underline{\underline{\mu}}, \quad (9)$$

$$M(\Phi)(\underline{\underline{\mu}}) = M_\Phi \times \underline{\underline{\mu}}, \quad (10)$$

$$M(\tilde{\Phi})(\underline{\underline{\mu}}, \underline{\underline{\lambda}}) = M_{\tilde{\Phi}} \times \underline{\underline{\mu}} \times \underline{\underline{\lambda}}. \quad (11)$$

We have

$$M_f \times \underline{\underline{\mu}} = M_f \cdot \underline{\underline{\mu}}, \quad (12)$$

$$M_\Phi \times \underline{\underline{\mu}} = M_\Phi \cdot \underline{\underline{\mu}}, \quad (13)$$

$$M_{\tilde{\Phi}} \times \underline{\underline{\mu}} \times \underline{\underline{\lambda}} = M_{\tilde{\Phi}} \cdot (\underline{\underline{\mu}} \times \underline{\underline{\lambda}}), \quad (14)$$

where ' \cdot ' is the product of the matrices.

Proof. We note first that

$$\mathbf{D} = L_{2 \times 1},$$

$$\mathbf{D}^{(n)} = L_{2^n \times 1}$$

are true. As far as $\underline{\underline{\mu}} \in L_{2^n \times 1}$ and $M_f \in L_{2 \times 2^n}$, we infer from Remark 2.3 that (12) holds. On the other hand, $\underline{\underline{\mu}} \in L_{2^n \times 1}$ makes $M_f \cdot \underline{\underline{\mu}}$ coincide with one of $\text{col}_1(M_f), \dots, \text{col}_{2^n}(M_f)$ and we know that $\text{col}_1(M_f), \dots, \text{col}_{2^n}(M_f) \in L_{2 \times 1}$, thus we can define $M(f)$ as

$$\mathbf{D}^{(n)} \ni \underline{\underline{\mu}} \mapsto M(f)(\underline{\underline{\mu}}) = M_f \times \underline{\underline{\mu}} \in \mathbf{D}.$$

The other statements are proved similarly. \square

Notation 4.3. We denote $F_{n,m} = \{h|h : \mathbf{B}^n \rightarrow \mathbf{B}^m\}$.

Theorem 4.4. (a) *The following diagrams*

$$\begin{array}{ccccc}
 \mathbf{B}^n & \xrightarrow{f} & \mathbf{B} & & \mathbf{B}^n & \xrightarrow{\Phi} & \mathbf{B}^n & & \mathbf{B}^n \times \mathbf{B}^n & \xrightarrow{\tilde{\Phi}} & \mathbf{B}^n \\
 \zeta_n \downarrow & & \downarrow \zeta & & \zeta_n \downarrow & & \downarrow \zeta_n & & \zeta_n \times \zeta_n \downarrow & & \downarrow \zeta_n \\
 \mathbf{D}^{(n)} & \xrightarrow{M(f)} & \mathbf{D} & & \mathbf{D}^{(n)} & \xrightarrow{M(\Phi)} & \mathbf{D}^{(n)} & & \mathbf{D}^{(n)} \times \mathbf{D}^{(n)} & \xrightarrow{M(\tilde{\Phi})} & \mathbf{D}^{(n)}
 \end{array}$$

commute.

(b) The assignments $F_{n,1} \ni f \mapsto M_f \in L_{2 \times 2^n}$, $F_{n,n} \ni \Phi \mapsto M_\Phi \in L_{2^n \times 2^n}$, $F_{2n,n} \ni \tilde{\Phi} \mapsto M_{\tilde{\Phi}} \in L_{2^n \times 2^{2n}}$ are bijective.

Proof. We fix $\mu \in \mathbf{B}^n$ and $\lambda \in \mathbf{B}^n$ arbitrary.

(a) In order to prove the commutativity of the first diagram, we use the fact that

$$\begin{aligned}
 f(\mu) &= f(1, \dots, 1, 1)\mu_1 \dots \mu_{n-1} \mu_n \oplus f(1, \dots, 1, 0)\mu_1 \dots \mu_{n-1} \overline{\mu_n} \\
 &\oplus f(1, \dots, 0, 1)\mu_1 \dots \overline{\mu_{n-1}} \mu_n \oplus \dots \oplus f(0, \dots, 0, 0)\overline{\mu_1} \dots \overline{\mu_{n-1}} \overline{\mu_n}, \\
 \overline{f(\mu)} &= \overline{f(1, \dots, 1, 1)\mu_1 \dots \mu_{n-1} \mu_n} \oplus \overline{f(1, \dots, 1, 0)\mu_1 \dots \mu_{n-1} \overline{\mu_n}} \\
 &\oplus \overline{f(1, \dots, 0, 1)\mu_1 \dots \overline{\mu_{n-1}} \mu_n} \oplus \dots \oplus \overline{f(0, \dots, 0, 0)\overline{\mu_1} \dots \overline{\mu_{n-1}} \overline{\mu_n}},
 \end{aligned}$$

wherefrom

$$\underline{\underline{f(\mu)}} = \left(\frac{f(\mu)}{\overline{f(\mu)}} \right) = M_f \cdot \underline{\underline{\mu}}. \quad (15)$$

We conclude that

$$(M(f) \circ \zeta_n)(\mu) = M(f)(\zeta_n(\mu)) = M(f)(\underline{\underline{\mu}}) \stackrel{(9)}{=} M_f \times \underline{\underline{\mu}}$$

$$\stackrel{(12)}{=} M_f \cdot \underline{\underline{\mu}} \stackrel{(15)}{=} \underline{\underline{f(\mu)}} = \zeta(f(\mu)) = (\zeta \circ f)(\mu),$$

i.e. the first diagram is commutative.

As far as the second diagram is concerned, we can prove that

$$\underline{\underline{\Phi(\mu)}} = M_\Phi \cdot \underline{\underline{\mu}}, \quad (16)$$

which is analogue with (15), and we obtain

$$(M(\Phi) \circ \zeta_n)(\mu) = M(\Phi)(\zeta_n(\mu)) = M(\Phi)(\underline{\underline{\mu}}) \stackrel{(10)}{=} M_\Phi \times \underline{\underline{\mu}}$$

$$\stackrel{(13)}{=} M_\Phi \cdot \underline{\underline{\mu}} \stackrel{(16)}{=} \underline{\underline{\Phi(\mu)}} = \zeta_n(\Phi(\mu)) = (\zeta_n \circ \Phi)(\mu).$$

For the commutativity of the third diagram, we have

$$\begin{aligned} \underline{\underline{\tilde{\Phi}(\mu, \lambda)}} &= \begin{pmatrix} \Phi_1^\lambda(\mu) \dots \Phi_{n-1}^\lambda(\mu) \overline{\Phi_n^\lambda(\mu)} \\ \Phi_1^\lambda(\mu) \dots \Phi_{n-1}^\lambda(\mu) \overline{\Phi_n^\lambda(\mu)} \\ \Phi_1^\lambda(\mu) \dots \Phi_{n-1}^\lambda(\mu) \overline{\Phi_n^\lambda(\mu)} \\ \dots \\ \Phi_1^\lambda(\mu) \dots \Phi_{n-1}^\lambda(\mu) \overline{\Phi_n^\lambda(\mu)} \end{pmatrix}, \\ \underline{\underline{\mu \times \lambda}} &= \begin{pmatrix} \mu_1 \dots \mu_n \lambda_1 \dots \lambda_{n-1} \lambda_n \\ \mu_1 \dots \mu_n \lambda_1 \dots \lambda_{n-1} \overline{\lambda_n} \\ \mu_1 \dots \mu_n \lambda_1 \dots \lambda_{n-1} \lambda_n \\ \dots \\ \overline{\mu_1} \dots \overline{\mu_n} \overline{\lambda_1} \dots \overline{\lambda_{n-1}} \overline{\lambda_n} \end{pmatrix} \end{aligned}$$

and we note that

$$\underline{\underline{\tilde{\Phi}(\mu, \lambda)}} = M_{\tilde{\Phi}} \cdot \underline{\underline{\mu \times \lambda}}. \quad (17)$$

For example, the second row in (17) is proved like this:

$$\begin{aligned} & \Phi_1^\lambda(\mu) \dots \Phi_{n-1}^\lambda(\mu) \overline{\Phi_n^\lambda(\mu)} \\ &= \Phi_1^{(1, \dots, 1, 1)}(1, \dots, 1) \dots \Phi_{n-1}^{(1, \dots, 1, 1)}(1, \dots, 1) \overline{\Phi_n^{(1, \dots, 1, 1)}(1, \dots, 1) \mu_1 \dots \mu_n \lambda_1 \dots \lambda_{n-1} \lambda_n} \\ & \oplus \Phi_1^{(1, \dots, 1, 0)}(1, \dots, 1) \dots \Phi_{n-1}^{(1, \dots, 1, 0)}(1, \dots, 1) \overline{\Phi_n^{(1, \dots, 1, 0)}(1, \dots, 1) \mu_1 \dots \mu_n \lambda_1 \dots \lambda_{n-1} \overline{\lambda_n}} \\ & \quad \dots \\ & \oplus \Phi_1^{(0, \dots, 0, 0)}(1, \dots, 1) \dots \Phi_{n-1}^{(0, \dots, 0, 0)}(1, \dots, 1) \overline{\Phi_n^{(0, \dots, 0, 0)}(1, \dots, 1) \mu_1 \dots \mu_n \overline{\lambda_1} \dots \overline{\lambda_{n-1}} \overline{\lambda_n}} \\ & \quad \dots \\ & \oplus \Phi_1^{(0, \dots, 0, 0)}(0, \dots, 0) \dots \Phi_{n-1}^{(0, \dots, 0, 0)}(0, \dots, 0) \overline{\Phi_n^{(0, \dots, 0, 0)}(0, \dots, 0) \overline{\mu_1} \dots \overline{\mu_n} \overline{\lambda_1} \dots \overline{\lambda_{n-1}} \overline{\lambda_n}}. \end{aligned}$$

We infer

$$\begin{aligned} (M(\tilde{\Phi}) \circ (\zeta_n \times \zeta_n))(\mu, \lambda) &= M(\tilde{\Phi})(\zeta_n(\mu), \zeta_n(\lambda)) = M(\tilde{\Phi})(\underline{\underline{\mu}}, \underline{\underline{\lambda}}) \stackrel{(11)}{=} M_{\tilde{\Phi}} \times \underline{\underline{\mu}} \times \underline{\underline{\lambda}} \\ &\stackrel{(14)}{=} M_{\tilde{\Phi}} \cdot \underline{\underline{\mu \times \lambda}} \stackrel{(17)}{=} \underline{\underline{\tilde{\Phi}(\mu, \lambda)}} = \zeta_n(\tilde{\Phi}(\mu, \lambda)) = (\zeta_n \circ \tilde{\Phi})(\mu, \lambda). \end{aligned}$$

(b) For example we suppose against all reason that $f, f' : \mathbf{B}^n \rightarrow \mathbf{B}$ exist, $f \neq f'$, with the property that $M_f = M_{f'}$. The hypothesis states the existence of $\mu \in \mathbf{B}^n$ such that $f(\mu) \neq f'(\mu)$ thus, from Theorem 3.4, $\underline{\underline{f(\mu)}} \neq \underline{\underline{f'(\mu)}}$. We have:

$$\begin{aligned} \underline{\underline{f(\mu)}} &= \zeta(f(\mu)) = (\zeta \circ f)(\mu) = (M(f) \circ \zeta_n)(\mu) = M(f)(\zeta_n(\mu)) \\ &= M(f)(\underline{\underline{\mu}}) = M_f \times \underline{\underline{\mu}} = M_{f'} \times \underline{\underline{\mu}} = M(f')(\underline{\underline{\mu}}) = M(f')(\zeta_n(\mu)) \\ &= (M(f') \circ \zeta_n)(\mu) = (\zeta \circ f')(\mu) = \zeta(f'(\mu)) = \underline{\underline{f'(\mu)}}, \end{aligned}$$

contradiction, showing that the assignment $F_{n,1} \ni f \mapsto M_f \in L_{2 \times 2^n}$ is injective. Due to the fact that $\text{card}(F_{n,1}) = \text{card}(L_{2 \times 2^n}) = 2^{2^n}$, injectivity and bijectivity coincide. \square

5. EQUATIONS OF EVOLUTION

Remark 5.1. Daizhan Cheng's theory adapted to asynchronicity replaces the equation of evolution (2) where $\Phi : \mathbf{B}^n \rightarrow \mathbf{B}^n$, $x : \mathbf{N} \rightarrow \mathbf{B}^n$, $\alpha \in \Pi_n$, $k \in \mathbf{N}$, with the equation

$$\underline{\underline{x(k+1)}} = \underline{\underline{\Phi^{\alpha^k}(x(k))}} = \underline{\underline{\widetilde{\Phi}(x(k), \alpha^k)}} \stackrel{(17)}{=} M_{\widetilde{\Phi}} \cdot \underline{\underline{x(k)}} \times \underline{\underline{\alpha^k}}, \quad (18)$$

which is easier to be studied. The price to pay is the increase of the dimension of the system from n to 2^n .

Example 5.1. We return to Example 1.1 now. Function

$$\widetilde{\Phi}(\mu_1, \mu_2, \lambda_1, \lambda_2) = (\overline{\lambda_1 \mu_1} \cup \lambda_1 \mu_2, \overline{\lambda_2 \mu_2} \cup \lambda_2 \overline{\mu_1})$$

defines the matrix

$$M_{\widetilde{\Phi}} = \begin{pmatrix} \overline{\widetilde{\Phi}_1(1, 1, 1, 1)} \overline{\widetilde{\Phi}_2(1, 1, 1, 1)} & \overline{\widetilde{\Phi}_1(0, 0, 0, 0)} \overline{\widetilde{\Phi}_2(0, 0, 0, 0)} \\ \overline{\widetilde{\Phi}_1(1, 1, 1, 1)} \overline{\widetilde{\Phi}_2(1, 1, 1, 1)} & \overline{\widetilde{\Phi}_1(0, 0, 0, 0)} \overline{\widetilde{\Phi}_2(0, 0, 0, 0)} \\ \overline{\widetilde{\Phi}_1(1, 1, 1, 1)} \overline{\widetilde{\Phi}_2(1, 1, 1, 1)} & \dots & \overline{\widetilde{\Phi}_1(0, 0, 0, 0)} \overline{\widetilde{\Phi}_2(0, 0, 0, 0)} \\ \overline{\widetilde{\Phi}_1(1, 1, 1, 1)} \overline{\widetilde{\Phi}_2(1, 1, 1, 1)} & \overline{\widetilde{\Phi}_1(0, 0, 0, 0)} \overline{\widetilde{\Phi}_2(0, 0, 0, 0)} \end{pmatrix}$$

$$= \begin{pmatrix} 0 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 \end{pmatrix}.$$

Equation (3) implies that $\forall k \in \mathbf{N}$, (18) is true, with $\underline{\underline{x, \alpha}} : \mathbf{N} \rightarrow \mathbf{D}^{(2)}$. We can see that, via (18), equation (4) becomes

$$\underline{\underline{x(0)}} \times \underline{\underline{\alpha^0}} = \underline{\underline{(0, 0)}} \times \underline{\underline{(1, 0)}} = \delta_4^4 \times \delta_4^2 = \delta_{16}^{14},$$

$$\underline{\underline{x(1)}} = M_{\widetilde{\Phi}} \cdot \delta_{16}^{14} = \delta_4^4 = \underline{\underline{(0, 0)}},$$

while (5) becomes

$$\underline{\underline{x(1)}} \times \underline{\underline{\alpha^1}} = \underline{\underline{(0, 0)}} \times \underline{\underline{(0, 1)}} = \delta_4^4 \times \delta_4^3 = \delta_{16}^{15},$$

$$\underline{\underline{x(2)}} = M_{\widetilde{\Phi}} \cdot \delta_{16}^{15} = \delta_4^3 = \underline{\underline{(0, 1)}},$$

(6) becomes

$$\underline{\underline{x(2)}} \times \underline{\underline{\alpha^2}} = \underline{\underline{(0, 1)}} \times \underline{\underline{(1, 1)}} = \delta_4^3 \times \delta_4^1 = \delta_{16}^9,$$

$$\underline{\underline{x(3)}} = M_{\widetilde{\Phi}} \cdot \delta_{16}^9 = \delta_4^1 = \underline{\underline{(1, 1)}},$$

(7) becomes

$$\begin{aligned} \underline{\underline{x(3)}} \times \underline{\underline{\alpha^3}} &= \underline{\underline{(1, 1)}} \times \underline{\underline{(0, 1)}} = \delta_4^1 \times \delta_4^3 = \delta_{16}^3, \\ \underline{\underline{x(4)}} &= M_{\overline{\Phi}} \cdot \delta_{16}^3 = \delta_4^2 = \underline{\underline{(1, 0)}}, \end{aligned}$$

and (8) becomes

$$\begin{aligned} \underline{\underline{x(4)}} \times \underline{\underline{\alpha^4}} &= \underline{\underline{(1, 0)}} \times \underline{\underline{(1, 0)}} = \delta_4^2 \times \delta_4^2 = \delta_{16}^6, \\ \underline{\underline{x(5)}} &= M_{\overline{\Phi}} \cdot \delta_{16}^6 = \delta_4^4 = \underline{\underline{(0, 0)}}, \\ &\dots \end{aligned}$$

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(Vlad Serban E.) ORADEA CITY HALL, DEPARTMENT OF COMPUTERS, ROMÂNIA
E-mail address: serban_e_vlad@yahoo.com